

07-1700

A

ASSISTANT COMMISSIONER OF PATENTS AND TRADEMARKS
Washington, DC 20231

PATENT
Date: July 14, 2000
File No. 0828.64472

07/14/00
U.S. PTO

Transmitted herewith for filing is the patent application of
Inventor(s): Masatoshi Haraguchi, Masakazu Hayashi and
Yuji Watanabe

I hereby certify that this paper is being deposited with the United States Postal Service as Express Mail in an envelope addressed to: Asst. Comm. for Patents, Washington, D.C. 20231, on this date.

7-14-00
Date

Express Mail Label No.: EL409491594US

For: METHOD AND APPARATUS FOR ANALYZING
PROGRAM BASED ON RULES OF PROGRAMMING
LANGUAGE BEFORE OPTIMIZATION IN COMPILER

Enclosed are:

- (X) 33 pages of specification, including 20 claims and an abstract.
- (X) an executed oath or declaration, with power of attorney.
- () an unexecuted oath or declaration, with power of attorney.
- () sheet(s) of informal drawing(s).
- (X) 9 sheet(s) of formal drawings(s).
- (X) Assignment(s) of the invention to FUJITSU LIMITED.
- (X) Assignment Form Cover Sheet.
- (X) A check in the amount of \$ 40.00 to cover the fee for recording the assignment(s) is enclosed.
- (X) Information Disclosure Statement.
- (X) Form PTO-1449 and cited references.
- () Associate power of attorney.
- (X) Priority Document.

Fee Calculation For Claims As Filed

a) Basic Fee						\$ 690.00
b) Independent Claims	<u>5</u>	-	<u>3</u>	=	<u>2</u>	x \$ 78.00 = \$ <u>156.00</u>
c) Total Claims	<u>20</u>	-	<u>20</u>	=	<u>0</u>	x \$ 18.00 = \$ <u> </u>
d) Fee for Multiple Claims						\$260.00 = \$ <u> </u>
Total Filing Fee						\$ <u>846.00</u>

- (X) A check in the amount of \$ 846.00 to cover the filing fee is enclosed.
- (X) The Commissioner is hereby authorized to charge any additional fees which may be required to this application under 37 C.F.R. §§1.16-1.17, or credit any overpayment, to Deposit Account No. 07-2069. Should no proper payment be enclosed herewith, as by a check being in the wrong amount, unsigned, post-dated, otherwise improper or informal or even entirely missing, the Commissioner is authorized to charge the unpaid amount to Deposit Account No. 07-2069. A duplicate copy of this sheet is enclosed.

Suite 8660 - Sears Tower
233 S. Wacker Drive
Chicago, Illinois 60606
(312) 993-0080

GREER, BURNS & CRAIN, LTD.

By:

Patrick G. Burns

Registration No. 29,367

U.S. PTO
09/616799
07/14/00

7/14/00
Date

Express Mail Label No. EL 409491594US

Signature



METHOD AND APPARATUS FOR ANALYZING PROGRAM BASED ON
RULES OF PROGRAMMING LANGUAGE BEFORE OPTIMIZATION IN
COMPILER

5 BACKGROUND OF THE INVENTION

1) Field of the Invention

The present invention relates to a method for facilitating optimization in a compiler. The present invention also relates to a compiler in which provision
10 is made for facilitating the optimization. The present invention further relates to an apparatus for facilitating optimization processing in a compiler.

2) Description of the Related Art

In the usual compilers, syntax analysis and the
15 like of a source code are performed at the front end, and an object code is generated at the back end. In addition, the object code is optimized for enhancing efficiency of execution of the object code. For example, when the program includes an array description, the
20 array may be recognized as a loop. Since loop operations occupy a considerable time in the execution of the object code, array portions of the object code is required to be optimized.

In some programming languages such as Fortran 90
25 and Fortran 95, values of or relationships between variables may be determined based on standardized rules or restrictions specified for the programming languages,

although such rules or restrictions are specified in different manners depending on the programming languages. For example, in Fortran 90 or Fortran 95, scalar representations of arrays as illustrated in Fig. 8(A) can be rewritten in vector representations as illustrated in Fig. 8(B), while in C and FORTRAN 77, access to arrays can be described in only element by element. According to the rules and restrictions of Fortran 90 or Fortran 95, the value of the variable L in the description in the fourth line of the text illustrated in Fig. 8(B) can be determined to be 100.

Loop fusion is one of the well-known optimization techniques, and is a technique of decreasing the number of loops for reducing overheads which are required for loop control. That is, a plurality of loops satisfying a predetermined condition are combined into a single loop. For loop fusion, some sophisticated techniques are known for analyzing relationships between arrays.

Further, Japanese Unexamined Patent Publication, No. 9-62514 discloses a method for facilitating loop fusion processing which is performed for optimizing the program code. In the method, a technique for analyzing relationships between arrays is used. To be more specific, JPP 9-62514 discloses facilitation of loop fusion processing by using a function of analyzing subscripts of the arrays in a program code, which is described in a programming language allowing vector

representations of arrays, such as Fortran 90 and Fortran 95. For example, Fig. 9 exhibits an example of a result of loop fusion, which is obtained by the subscript analysis of arrays in a program code. The

5 subscript analysis is performed on the three assignment statements expressed by arrays, which are indicated in the upper rectangle in Fig. 9. The result of the subscript analysis is utilized for producing an optimum (fused) loop, as indicated in the lower rectangle in Fig.

10 9. According to the rules and restrictions specified for the above programming language such as Fortran 90 and Fortran 95, the sizes of the arrays on the right-hand side and the left-hand side of each assignment statement are identical. Therefore, the result of subscript

15 analysis performed on the arrays in the three assignment statements in the upper rectangle in Fig. 9 shows that $L=5=M-1=N-2$. This result is used for facilitation of the loop fusion processing and production of a fused loop having an optimum form, as illustrated in Fig. 9. For

20 example, the number of iterations in the fused DO loop is determined as $L=5$ by the above subscript analysis, in advance of the loop fusion processing.

As described above, conventionally, the result of subscript analysis of arrays is utilized for

25 facilitating the loop fusion processing. However, the result of subscript analysis of arrays is not utilized in optimization processing other than the loop fusion

processing.

SUMMARY OF THE INVENTION

An object of the present invention is to provide a
5 method for facilitating various types of optimization in
a compiler, based on rules or restrictions specified for
a programming language of a program code.

Another object of the present invention is to
provide a compiler in which provision is made for
10 facilitating various types of optimization, based on
rules or restrictions specified for a programming
language of a program code.

A further object of the present invention is to
provide an apparatus for facilitating optimization
15 processing in a compiler, whereby various types of
optimization can be facilitated based on rules or
restrictions specified for a programming language of a
program code.

(1) According to the first aspect of the present
20 invention, there is provided a method of facilitating
optimization processing in a compiler. The method
comprises the steps of: (a) storing, in a language-
specific-rule table, one or more predetermined rules
which are specified for one or more programming
25 languages; (b) analyzing a program code which includes
one or more instructions, and is described in one of the
one or more programming languages, based on the one or

more predetermined rules, to obtain an analysis result;
and (c) embedding the analysis result in the program
code.

5 The method according to the first aspect of the
present invention may also have one or any possible
combination of the following additional features (i) to
(v).

(i) In the step (a), the one or more
predetermined rules may be stored in the language-
10 specific-rule table as one or more language-specific-
information analyzing functions, and the step (b) may
comprise the substeps of (d) reading out, from the
language-specific-rule table, at least one of the one or
more language-specific-information analyzing functions
15 which is needed for analyzing the program code, and (e)
determining values of or relationships between variables
included in the program code, based on the at least one
of the one or more language-specific-information
analyzing functions read out in the step (d), and
20 producing the analysis result which includes the
determined values of or relationships between the
variables.

(ii) In addition to the feature (i), the
step (b) may further comprise the substep of (f) the at
25 least one of the one or more language-specific-
information analyzing functions read out in the step (d)
is registered in a check function table for use in the

step (e).

(iii) In addition to the feature (ii), the operation in step (b) may be performed for each instruction set which is comprised of at least one of the one or more instructions, and in the step (c), the analysis result may be embedded in a position preceding the instruction set in the program code.

(iv) The program code may be a source code.

(v) The program code may be an intermediate code.

(2) According to the second aspect of the present invention, there is provided an apparatus for facilitating optimization processing in a compiler. The apparatus comprises a language-specific-rule table which stores one or more predetermined rules which are specified for one or more programming languages; an analyzing unit which analyzes a program code which includes one or more instructions, and is described in one of the one or more programming languages, based on the one or more predetermined rules, to obtain an analysis result; and an embedding unit which embeds the analysis result in the program code.

The apparatus according to the second aspect of the present invention may also have one or any possible combination of the following additional features (vi) to (x).

(vi) The language-specific-rule table may

store the one or more predetermined rules as one or more language-specific-information analyzing functions, and the analyzing unit may comprise a readout unit which reads out, from the language-specific-rule table, at least one of the one or more language-specific-information analyzing functions which is needed for analyzing the program code, and a determination unit which determines values of or relationships between variables included in the program code, based on the at least one of the one or more language-specific-information analyzing functions read out by the readout unit, and produces the analysis result which includes the determined values of or relationships between the variables.

(vii) In addition to the feature (vi), the analyzing unit may comprise a check function table in which the at least one of the one or more language-specific-information analyzing functions read out by the readout unit is registered for use by the determination unit.

(viii) In addition to the feature (vii), the operation of the analyzing unit may be performed for each instruction set which is comprised of at least one of the one or more instructions, and the embedding unit may embed the analysis result in a position preceding the instruction set in the program code.

(ix) The program code may be a source code.

(x) The program code may be an intermediate code.

(3) According to the third aspect of the present invention, there is provided a product for use with an apparatus for facilitating optimization processing in a compiler. When the product is used with the apparatus, the product is able to output control information which directs the apparatus to comprise: a language-specific-rule table which stores one or more predetermined rules which are specified for one or more programming languages; an analyzing unit which analyzes a program code which includes one or more instructions, and is described in one of the one or more programming languages, based on the one or more predetermined rules, to obtain an analysis result; and an embedding unit which embeds the analysis result in the program code.

The product according to the third aspect of the present invention may have one or any possible combination of the aforementioned additional features (vi) to (x).

(4) According to the fourth aspect of the present invention, there is provided a compiler comprising: a front end unit which performs syntax analysis of a source code which is described in one of one or more predetermined programming languages, to produce an intermediate code; a language-specific-rule table which stores one or more predetermined rules which are

specified for the one or more predetermined programming languages; an analyzing unit which analyzes the intermediate code which includes one or more instructions, based on the one or more predetermined rules, to obtain an analysis result; an embedding unit which embeds the analysis result in the program code to produce a modified intermediate code; and an optimizing unit which performs an operation of optimizing the modified intermediate code.

10 The compiler according to the fourth aspect of the present invention may have one or any possible combination of the aforementioned additional features (vi) to (viii).

(5) According to the fifth aspect of the present invention, there is provided a compiler comprising: a language-specific-rule table which stores one or more predetermined rules which are specified for one or more programming languages; an analyzing unit which analyzes a source code which includes one or more instructions, and is described in one of the one or more programming languages, based on the one or more predetermined rules, to obtain an analysis result; an embedding unit which embeds the analysis result in the source code to produce a modified source code; a syntax analyzing unit which performs syntax analysis of the modified source code which is described in one of one or more predetermined programming languages, to produce an intermediate code;

and an optimizing unit which performs an optimization operation on the intermediate code.

The compiler according to the fifth aspect of the present invention may have one or any possible
5 combination of the aforementioned additional features (vi) to (viii).

(6) According to the first to fifth aspects of the present invention, the program code is analyzed based on rules or restrictions which are specified for the
10 programming language of the program code so as to determine values of or relationships between variables included in the program code, and information on the obtained values of or relationships between variables is embedded in the program code. Therefore, the
15 optimization function of the compiler can easily utilize the analysis result when the optimization function handle the program code. Thus, the optimization processing is facilitated by the provision of the present invention, and the type of the facilitated
20 optimization processing is not limited to the loop fusion.

The above and other objects, features and advantages of the present invention will become apparent from the following description when taken in conjunction
25 with the accompanying drawings which illustrate preferred embodiment of the present invention by way of example.

BRIEF DESCRIPTION OF THE DRAWINGS

In the drawings:

Fig. 1 is a diagram illustrating the construction
5 of a compiler as an embodiment of the fourth aspect of
the present invention;

Figs. 2(A) and 2(B) are diagrams illustrating
examples of contents of the language-specific-rule table
7;

10 Fig. 3 is a diagram illustrating an example of a
source code;

Fig. 4 is a diagram illustrating a modified source
code which is produced by embedding the analysis results,
L=5 and N=6, in the source code of Fig. 3;

15 Fig. 5 is a diagram illustrating a result of
optimization obtained by utilizing the analysis result,
L=5, and completely unrolling the loops included in the
source code of Fig. 3;

Fig. 6 is a diagram illustrating a further
20 optimized program code obtained by enhancing the degree
of parallelism in the program code of Fig. 5;

Fig. 7 is a flowchart of the operation of the
optimization facilitation unit 5 in Fig. 1;

Fig. 8(A) is a diagram of an example of a program
25 code including scalar representations of arrays;

Fig. 8(B) is a diagram of an example of a program
code including vector representations of arrays; and

Fig. 9 exhibits an example of a result of loop fusion, which is obtained by the subscript analysis of arrays in a program code.

5 DESCRIPTION OF THE PREFERRED EMBODIMENTS

Embodiments of the present invention are explained below with reference to drawings.

(1) Construction of Compiler

Fig. 1 is a diagram illustrating the construction of a compiler as an embodiment of the fourth aspect of the present invention. The compiler of Fig. 1 comprises an input unit 1, a front end unit 2, an optimization unit 3, an output unit 4, and an optimization facilitation unit 5. The input unit 1 inputs a source program (code). The front end unit 2 performs syntax analysis and the like of the source code, and produces an intermediate code. The optimization unit 3 processes the intermediate code so as to optimize an object program (code). The output unit 4 produces the object code based on the intermediate code processed by the optimization unit 3.

The optimization facilitation unit 5 is an embodiment of the apparatus according to the second aspect of the present invention, and is provided between the front end unit 2 and the optimization unit 3 in the embodiment of Fig. 1. The optimization facilitation unit 5 analyzes the intermediate code based on the rules and

restrictions specified for one of a plurality of programming languages, and modifies the intermediate code so as to facilitate the optimization processing which is performed by the optimization unit 3. The optimization facilitation unit 5 comprises a language-specific-information analyzing unit 6, a language-specific-rule table 7, an analysis result embedding unit 8, and a check function table 9. The language-specific-rule table 7 stores information on the rules and restrictions which are specified for a plurality of programming languages, in the form of one or more language-specific-information analyzing functions (hereinafter also called check functions).

When a source code is input into the input unit 1, the front end unit 2 performs syntax analysis and the like of the source code, and produces an intermediate code. In the optimization facilitation unit 5, the language-specific-information analyzing unit 6 first refers to the language-specific-rule table 7 to determine at least one of the one or more check functions which is used for analyzing language specific information in the intermediate code, and registers the determined check function in the check function table 9, which is provided in the analysis result embedding unit 8. Next, when the language-specific-information analyzing unit 6 receives each instruction in the intermediate code produced by the front end unit 2, the

language-specific-information analyzing unit 6
determines whether or not values of or relationships
between the variables in the instruction can be
determined, and obtains the values of or relationships
5 between the variables when they can be determined. Thus,
an analysis result is obtained for each single
instruction in the intermediate code, or for each set of
a plurality of instructions in the intermediate code,
when it is possible. Then, the analysis result embedding
10 unit 8 embeds the obtained analysis result in a position
preceding each instruction in the intermediate code,
where the obtained analysis result is described in the
form which can be recognized by the optimization unit 3.
Thus, the intermediate code is modified by the analysis
15 result embedding unit 8. The modified intermediate code
is passed to the optimization unit 3. In the
optimization unit 3, the intermediate code is processed
so as to optimize an object code, which is then produced
by the output unit 4 based on the optimized intermediate
20 code produced by the optimization unit 3.

As described above, the analysis result is
described in the form which can be recognized by the
optimization unit 3, and embedded in the intermediate
code, where the analysis result is obtained
25 corresponding to each instruction in the intermediate
code, and indicates the values of or relationships
between the variables in the intermediate code.

Therefore, for example, the optimization unit 3 can receive a constant as a value of a variable, instead of receiving as the variable. Thus, the optimization unit 3 can utilize the above analysis result for not only the
5 loop fusion but also all the other types of optimization.

In the construction of Fig. 1, the optimization facilitation unit 5 is arranged between the front end unit 2 and the optimization unit 3, and the optimization facilitation unit 5 analyzes the intermediate code
10 produced by the front end unit 2, and embeds the analysis result in the intermediate code. However, the optimization facilitation unit 5 may be provided in the front end unit 2 so that the optimization facilitation unit 5 analyzes the source code before the syntax
15 analysis, and embeds the result of the analysis based on the rules and restrictions specified the programming language, in the source code. In this case, the compiler realizes the aforementioned fifth aspect of the present invention.

20 (2) Language-Specific-Rule Table

Figs. 2(A) and 2(B) are diagrams illustrating examples of contents of the language-specific-rule table 7. The language-specific-rule table 7 in this embodiment includes a first and second tables, as shown in Figs.
25 2(A) and 2(B), respectively. For example, the first table includes, as a language-specific-information analyzing function (check function), information

indicating whether or not the respective programming languages allow vector representation of arrays, as illustrated in Fig. 2(A). The second table includes at least one other language-specific-information analyzing function (at least one other check function), and indicates whether or not a certain rule corresponding to each of the at least one other check function is specified for the respective programming languages. For example, the rule (check function) "No. 1" illustrated in Fig. 2(B) is that the sizes of arrays on the right-hand and left-hand sides of each assignment statement are identical.

(3) Examples

Fig. 3 is a diagram illustrating an example of a source code. When the source code of Fig. 3 is written in Fortran 90 or Fortran 95, and the language-specific-information analyzing unit 6 analyzes the assignment statement in the fourth line of the source code of Fig. 3, $L=5$ is obtained as an analysis result based on the above rule No. 1 in Fig. 2(B). Similarly, $N=6$ is obtained as an analysis result from the assignment statement in the fifth line of the source code of Fig. 3, based on the above rule No. 1. According to the present invention, the above analysis results, $L=5$ and $N=6$, are passed to the optimization unit 3 by embedding the analysis results in the source code or an intermediate code by the analysis result embedding unit 8, so as to

facilitate optimization processing executed by the optimization unit 3.

Fig. 4 is a diagram illustrating a modified source code which is produced by embedding the above analysis results, $L=5$ and $N=6$, in the source code of Fig. 3. In the example of Fig. 4, the analysis result, $L=5$, is embedded in a position preceding the assignment statement which includes the variable L , and the other analysis result, $N=6$, is embedded in a position preceding the assignment statement which includes the variable N . However, in this example, only the former analysis result, $L=5$, facilitates the optimization processing by the optimization unit 3, as explained below.

Fig. 5 is a diagram illustrating a result of optimization obtained by utilizing the above analysis result, $L=5$, and completely unrolling the loops included in the source code of Fig. 3. In the program code of Fig. 5, the substantial loops represented by the assignment statements in the fourth and fifth lines, and the DO loop in the sixth to eighth lines of the source code of Fig. 3 are each unrolled completely in the program code of Fig. 5, and the above analysis result, $L=5$, is utilized when the DO loop is unrolled. That is, the loop unrolling optimization is facilitated by the analysis result, $L=5$, which is supplied to the optimization unit 3 by embedding the analysis result in a program code

supplied to the optimization unit 3.

Fig. 6 is a diagram illustrating a further optimized program code obtained by enhancing the degree of parallelism in the program code of Fig. 5, where the degree of parallelism is enhanced by changing the order of the assignment statements of Fig. 5. Generally, instructions, e.g., assignment statements, included in a program code must be executed in the order in which the instructions are written in the program code. That is, each instruction (assignment statement) cannot be executed unless the preceding instruction (assignment statement) is executed. However, in the program code of Fig. 6, the assignment operation, $k1=a(2)+a(3)$, is not affected by the result of the preceding assignment operation, $k=k+a(1)$. Therefore, the assignment operation, $k1=a(2)+a(3)$, can be executed without waiting for completion of the preceding assignment, $k=k+a(1)$. Similarly, the assignment operation, $k2=a(4)+a(5)$, can be executed without waiting for completion of the preceding assignment operations, $k=k+a(1)$ and $k1=a(2)+a(3)$. Thus, when a computer can perform a plurality of add operations in parallel, the above assignment operations, $k=k+a(1)$, $k1=a(2)+a(3)$, and $k2=a(4)+a(5)$, can be concurrently executed, although an error may be produced in the case where the arrays A and the variable k are represented by floating point numbers.

As explained above, the embedment of the analysis

result in the program code realizes constant propagation to the number of iterations in the DO loop, and thus facilitates the optimization processing.

Although the above explanations are provided based on the source (human-readable) programs as illustrated in Figs. 3 to 6, the program code, which is analyzed, and in which the analysis result is embedded, may be a source code, an intermediate code, or a machine-executable code such as the object code.

10 (4) Operation Flow

Fig. 7 is a flowchart of the operation of the optimization facilitation unit 5 in Fig. 1 which is an embodiment of the apparatus for facilitating optimization processing in a compiler, according to the second aspect of the present invention.

In step S1, the language-specific-information analyzing unit 6 refers to the language-specific-rule table 7, and determines at least one language-specific-information analyzing function (check function) which is used for analyzing language specific information in the intermediate code, and registers the determined check function in the check function table 9, which is provided in the analysis result embedding unit 8. For example, the language-specific-information analyzing unit 6 refers to the aforementioned first and second tables, determines the rules (check functions) Nos. 1 and 2 in the second table as the check functions to be

used in the analysis of the intermediate code, and registers the determined check functions in the check function table 9.

In step S2, the language-specific-information
5 analyzing unit 6 inputs an instruction from the intermediate code.

In step S3, it is determined whether or not any instruction remains in the intermediate code. When yes is determined, the operation goes to step S4. When no is
10 determined, the operation of Fig. 7 is completed.

In step S4, the count "i" is set to its initial value, "1".

In step S5, it is determined whether or not the count "i" exceeds the size of the check function table 9.
15 The size of the check function table 9 is the number of check functions registered in the check function table 9. For example, when only the rules (check functions) Nos. 1 and 2 are registered in the check function table 9, the above number is two. When yes is determined in step
20 S5, the operation goes to step S2. When no is determined in step S5, the operation goes to step S6.

In step S6, the language-specific-information analyzing unit 6 calls the i-th check function from the check function table 9.

25 In step S7, the analysis result embedding unit 8 analyzes the instruction by using the called check function, to obtain an analysis result. For example,

"L=5" and "N=6" are obtained as the analysis results by analyzing the instructions in the fourth and fifth lines of the source code of Fig. 3.

In step S8, the analysis result embedding unit 8 embeds the obtained analysis result in a position preceding the analyzed instruction in the program.

In step S9, the count "i" is incremented, and the operation goes to step S5.

As explained above, the analysis is performed by each of the check functions to obtain an analysis result, and the analysis result is embedded in the program. Thus, for example, the intermediate code as illustrated in Fig. 4 is obtained. Then, the program in which the analysis result is embedded is input into the optimization unit 3, and therefore the optimization as illustrated in Figs. 5 and 6 can be performed.

(5) Other Matters

(i) Instead of embedding the analysis result in a program code, all of the analysis results may be held in a table so that the optimization unit 3 can utilize the contents of the table. However, according to this method, bothersome analyzing operations are required for the optimization unit 3 to obtain information on values or effective ranges of variables. In addition, the optimization unit 3 must determine whether or not a value of a variable which is referred to in each instruction is available, by searching for the analysis

result in the table, while the optimization unit 3 can immediately use the analysis result in the usual optimization processing when the analysis result is embedded in the program code to be optimized.

5 (ii) In addition, the operations and the functions of the present invention may be realized by using a certain product with an apparatus provided for facilitating the optimization processing in a compiler, e.g., by installing a computer-readable medium in a
10 computer. The product is such that when the product is used with the apparatus (e.g., a computer), the product is able to output control information which directs the apparatus to realize any of the functions of the present invention. The product may be a semiconductor storage
15 device storing a program which realizes the above functions, such as a ROM, or a magnetic storage medium such as a floppy disc or a hard disk, or a CD-ROM, a CD-R, a DVD-ROM, a DVD-RAM, a DVD-R, or the like. Further, the above product may be a programmed hardware logic
20 circuit such as an LSI. The above product per se can be put into the market. Alternatively, program data realizing the above functions may be transferred through a communication network from a storage device included in a computer system to another computer system. When
25 the program is executed in a computer system, for example, the program stored in a hard disk may be loaded in a main memory of the computer system.

(iii) The foregoing is considered as illustrative only of the principle of the present invention. Further, since numerous modifications and changes will readily occur to those skilled in the art, it is not desired to
5 limit the invention to the exact construction and applications shown and described, and accordingly, all suitable modifications and equivalents may be regarded as falling within the scope of the invention in the appended claims and their equivalents.

10 (iv) All of the contents of the Japanese patent application, No. 11-305026 are incorporated into this specification by reference.

What is claimed is:

1. A method of facilitating optimization processing in a compiler, comprising the steps of:

(a) storing, in a language-specific-rule table, one or more predetermined rules which are specified for one or more programming languages;

(b) analyzing a program code which includes one or more instructions, and is described in one of said one or more programming languages, based on said one or more predetermined rules, to obtain an analysis result; and

(c) embedding said analysis result in said program code.

2. A method according to claim 1, wherein in said step (a), said one or more predetermined rules are stored in said language-specific-rule table as one or more language-specific-information analyzing functions, and

said step (b) comprises the substeps of,

(d) reading out, from said language-specific-rule table, at least one of said one or more language-specific-information analyzing functions which is needed for analyzing said program code, and

(e) determining values of or relationships between variables included in said program code, based on said at least one of said one or more language-

specific-information analyzing functions read out in said step (d), and producing said analysis result which includes the determined values of or relationships between the variables.

5

3. A method according to claim 2, wherein said step (b) further comprises the substep of,

(f) said at least one of said one or more language-specific-information analyzing functions read out in said step (d) is registered in a check function table for use in said step (e).

4. A method according to claim 3, wherein the operation in step (b) is performed for each instruction set which is comprised of at least one of said one or more instructions, and

in said step (c), said analysis result is embedded in a position preceding said each instruction set in said program code.

20

5. A method according to claim 1, wherein said program code is a source code.

6. A method according to claim 1, wherein said program code is an intermediate code.

7. An apparatus for facilitating optimization

processing in a compiler, comprising:

a language-specific-rule table which stores one or more predetermined rules which are specified for one or more programming languages;

5 an analyzing unit which analyzes a program code which includes one or more instructions, and is described in one of said one or more programming languages, based on said one or more predetermined rules, to obtain an analysis result; and

10 an embedding unit which embeds said analysis result in said program code.

8. An apparatus according to claim 7, wherein said language-specific-rule table stores said one or
15 more predetermined rules as one or more language-specific-information analyzing functions, and

said analyzing unit comprises,

a readout unit which reads out, from said language-specific-rule table, at least one of said one
20 or more language-specific-information analyzing functions which is needed for analyzing said program code, and

a determination unit which determines values of or relationships between variables included in said
25 program code, based on said at least one of said one or more language-specific-information analyzing functions read out by said readout unit, and produces said

analysis result which includes the determined values of or relationships between the variables.

9. An apparatus according to claim 8, wherein
5 said analyzing unit comprises,

a check function table in which said at least one of said one or more language-specific-information analyzing functions read out by said readout unit is registered for use by said determination unit.

10

10. An apparatus according to claim 9, wherein the operation of said analyzing unit is performed for each instruction set which is comprised of at least one of said one or more instructions, and

15

said embedding unit embeds said analysis result in a position preceding said each instruction set in said program code.

11. An apparatus according to claim 7, wherein
20 said program code is a source code.

12. An apparatus according to claim 7, wherein said program code is an intermediate code.

25 13. A product for use with an apparatus for facilitating optimization processing in a compiler, said product, when used with said apparatus, is able to

output control information which directs the apparatus to comprise:

a language-specific-rule table which stores one or more predetermined rules which are specified for
5 one or more programming languages;

an analyzing unit which analyzes a program code which includes one or more instructions, and is described in one of said one or more programming languages, based on said one or more predetermined rules,
10 to obtain an analysis result; and

an embedding unit which embeds said analysis result in said program code.

14. A product according to claim 13, wherein said
15 language-specific-rule table stores said one or more predetermined rules as one or more language-specific-information analyzing functions, and

said analyzing unit comprises,

a readout unit which reads out, from said
20 language-specific-rule table, at least one of said one or more language-specific-information analyzing functions which is needed for analyzing said program code, and

a determination unit which determines values
25 of or relationships between variables included in said program code, based on said at least one of said one or more language-specific-information analyzing functions

read out by said readout unit, and produces said analysis result which includes the determined values of or relationships between the variables.

5 15. A product according to claim 13, wherein said program code is a source code.

16. A product according to claim 13, wherein said program code is an intermediate code.

10

17. A compiler comprising:

15 a front end unit which performs syntax analysis of a source code which is described in one of one or more predetermined programming languages, to produce an intermediate code;

 a language-specific-rule table which stores one or more predetermined rules which are specified for said one or more predetermined programming languages;

20 an analyzing unit which analyzes said intermediate code which includes one or more instructions, based on said one or more predetermined rules, to obtain an analysis result;

25 an embedding unit which embeds said analysis result in said program code to produce a modified intermediate code; and

 an optimizing unit which performs an operation of optimizing said modified intermediate code.

18. A compiler according to claim 17, wherein
said language-specific-rule table stores said one or
more predetermined rules as one or more language-
5 specific-information analyzing functions, and

said analyzing unit comprises,

a readout unit which reads out, from said
language-specific-rule table, at least one of said one
or more language-specific-information analyzing
10 functions which is needed for analyzing said
intermediate code, and

a determination unit which determines values
of or relationships between variables included in said
intermediate code, based on said at least one of said
15 one or more language-specific-information analyzing
functions read out by said readout unit, and produces
said analysis result which includes the determined
values of or relationships between the variables.

20 19. A compiler comprising:

a language-specific-rule table which stores
one or more predetermined rules which are specified for
one or more programming languages;

an analyzing unit which analyzes a source
25 code which includes one or more instructions, and is
described in one of said one or more programming
languages, based on said one or more predetermined rules,

to obtain an analysis result;

an embedding unit which embeds said analysis result in said source code to produce a modified source code;

5 a syntax analyzing unit which performs syntax analysis of said modified source code which is described in one of one or more predetermined programming languages, to produce an intermediate code; and

10 an optimizing unit which performs an optimization operation on said intermediate code.

20. A compiler according to claim 19, wherein said language-specific-rule table stores said one or
15 more predetermined rules as one or more language-specific-information analyzing functions, and

said analyzing unit comprises,

a readout unit which reads out, from said language-specific-rule table, at least one of said one
20 or more language-specific-information analyzing functions which is needed for analyzing said source code, and

a determination unit which determines values of or relationships between variables included in said
25 source code, based on said at least one of said one or more language-specific-information analyzing functions read out by said readout unit, and produces said

analysis result which includes the determined values of
or relationships between the variables.

ABSTRACT OF THE DISCLOSURE

An apparatus for facilitating optimization
5 processing in a compiler includes a language-specific-
rule table which stores one or more predetermined rules
which are specified for one or more programming
languages; an analyzing unit which analyzes a program
code which includes one or more instructions, and is
10 described in one of the one or more programming
languages, based on the one or more predetermined rules,
to obtain an analysis result; and an embedding unit
which embeds the analysis result in the program code.

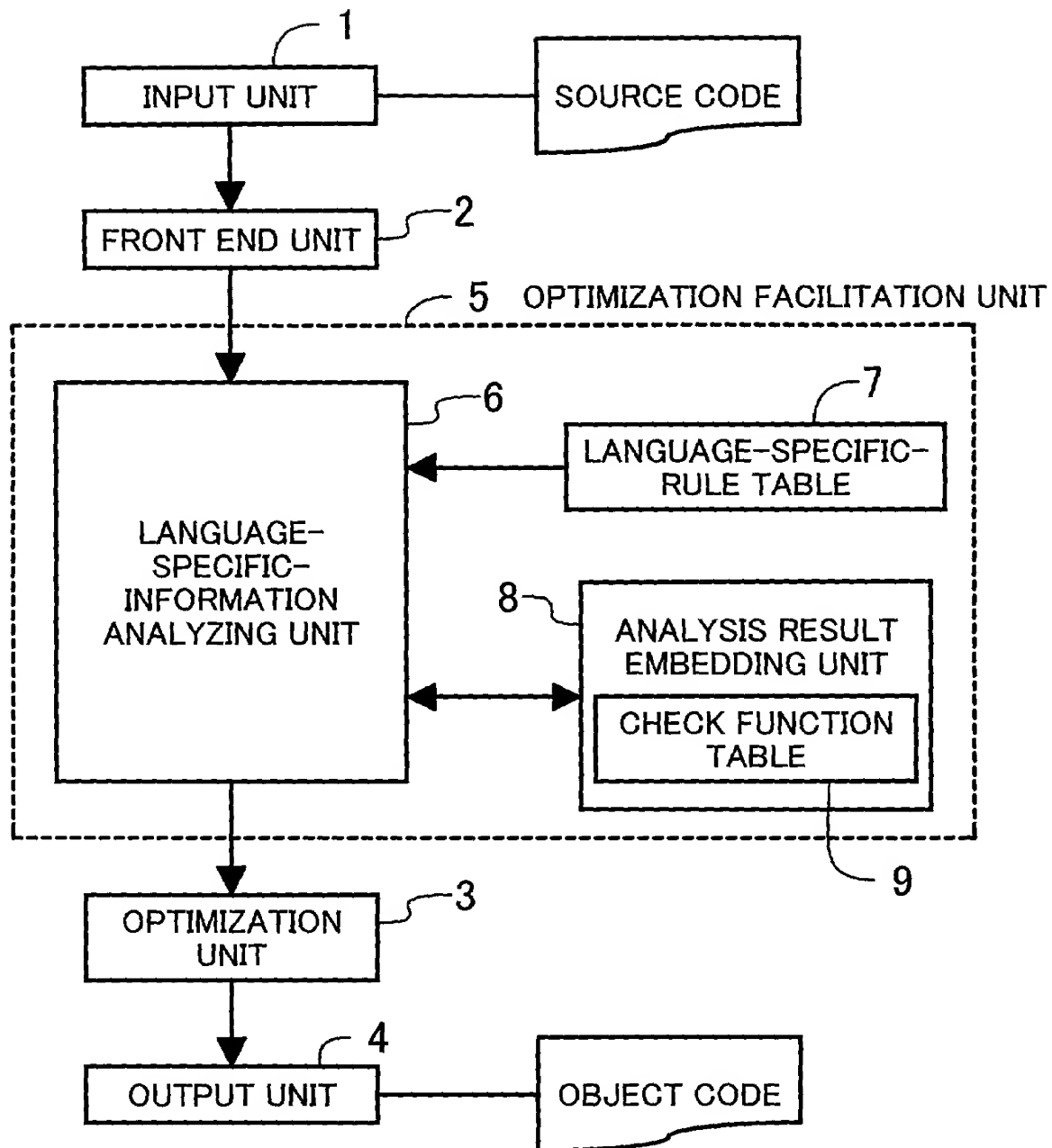


FIG. 1

FIG. 2(A)

LANGUAGE	Fortran77	Fortran90	Fortran95	C	Java	□□	△△
VECTOR REPRESENTATION OF ARRAYS	x	○	○	x	x	○	x
:	:	:	:	:	:	:	:

FIG. 2(B)

No.	LANGUAGE	Fortran90	Fortran95	□□
1	The sizes of arrays on the right-hand and left-hand sides of each assignment statement are identical.	○	○	x
2	The mask size of an array in a where statement is identical with the size of an array on the right-hand of side of each assignment statement in the where statement.	○	○	x
:	:	:	:	:

EXAMPLE OF SOURCE CODE

```
1 subroutine sub(a,b,k,L,N)
2 integer(kind=4), dimension(1:6) :: a,b
3 integer(kind=4) :: L,N
4 a(2:L) = b(2:5)
5 b(3:N) = a(1:4)
6 do i=1,L
7   k = k + a(i)
8 enddo
9 end subroutine
```

FIG. 3

EMBEDMENT OF ANALYSIS RESULT

```
subroutine sub(a,b,k,L,N)
integer(kind=4), dimension(1:6) :: a,b
integer(kind=4) :: L,N
L=5
a(2:L) = b(2:5)
N=6
b(3:N) = a(1:4)
do i=1,L
  k = k + a(i)
enddo
```

FIG. 4

RESULT OF COMPLETE UNROLLING OF LOOPS

$$\begin{aligned} a(2) &= b(2) \\ a(3) &= b(3) \\ a(4) &= b(4) \\ a(5) &= b(5) \\ b(3) &= a(1) \\ b(4) &= a(2) \\ b(5) &= a(3) \\ b(6) &= a(4) \\ k &= k + a(1) \\ k &= k + a(2) \\ k &= k + a(3) \\ k &= k + a(4) \\ k &= k + a(5) \end{aligned}$$

FIG. 5

RESULT OF OPTIMIZATION FOR
ENHANCING DEGREE OF PARALLELISM

```
a(2) = b(2)
a(3) = b(3)
a(4) = b(4)
a(5) = b(5)
b(3) = a(1)
b(4) = a(2)
b(5) = a(3)
b(6) = a(4)
k = k + a(1)
k1 = a(2) + a(3)
k2 = a(4) + a(5)
k = k + k1
k = k + k2
```

FIG. 6

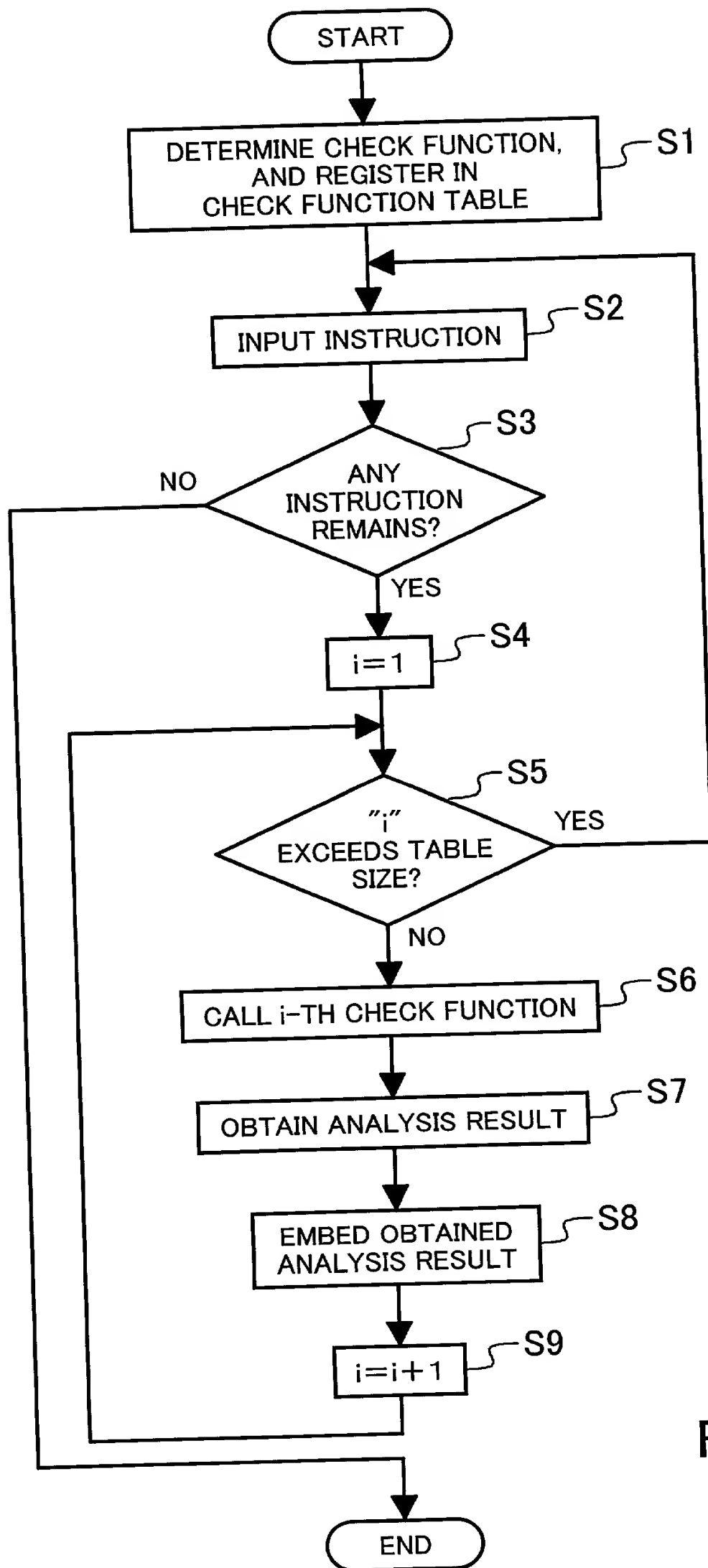


FIG. 7

SOURCE CODE INCLUDING SCALAR
REPRESENTATION OF ARRAYS

```
1 subroutine sub(r,q)
2 real r(100),q(100)
3 do i=1,100
4 r(i)=q(i)
5 enddo
6 end subroutine
```

PRIOR ART
FIG. 8(A)

SOURCE CODE INCLUDING VECTOR
REPRESENTATION OF ARRAYS

```
1 subroutine sub(r,q,L)
2 real r(100),q(100)
3 integer L
4 r(1:L)=q(1:100)
5 end subroutine
```

PRIOR ART
FIG. 8(B)

```
a(1:L)=b(2:6)
b(2:M)=c(3:L+2)
c(3:N)=a(1:M-1)
```

↓ L=5=M-1=N-2(RESET OF SUBSCRIPT
ANALYSIS OF ARRAYS)

```
do i=1,5 ! OPTIMIZED LOOP
  a(i)=b(i+1)
  b(i+1)=c(i+2)
  c(i+2)=a(i)
enddo
```

PRIOR ART
FIG. 9

Declaration and Power of Attorney For Patent Application

特許出願宣言書

Japanese Language Declaration.

私は、下欄に氏名を記載した発明者として、以下のとおり宣言する：

私の住所、郵便の宛先および国籍は、下欄に氏名に続いて記載したとおりであり、

名称の発明に関し、請求の範囲に記載した特許を求める主題の本来の、最初にして唯一の発明者である（一人の氏名のみが下欄に記載されている場合）か、もしくは本来の、最初にして共同の発明者である（複数の氏名が下欄に記載されている場合）と信じ、

その明細書を
(該当する方に印を付す)

☐ ここに添付する。

☐ _____ 日に出版番号
第 0 / _____ 号として提出し、

_____ 日に補正した。
(該当する場合)

私は、前記のとおり補正した請求の範囲を含む前記明細書の内容を検討し、理解したことを陳述する。

私は、連邦規則法典第37部第1章第56条(a)項に従い、本願の審査に所要の情報を開示すべき義務を有することを認める。

As a below named inventor, I hereby declare that:

My residence, post office address and citizenship are as stated below next to my name.

I believe I am the original, first and sole inventor (if only one name is listed below) or an original, first and joint inventor (if plural names are listed below) of the subject matter which is claimed and for which a patent is sought on the invention entitled

METHOD AND APPARATUS FOR ANALYZING

PROGRAM BASED ON RULES OF PROGRAMMING

LANGUAGE BEFORE OPTIMIZATION IN COMPILER

the specification of which

(check one)

☒ is attached hereto.

☐ was filed on _____ as

Application Serial No. 0 / _____

and was amended on _____
(if applicable)

I hereby state that I have reviewed and understand the contents of the above identified specification, including the claims, as amended by any amendment referred to above.

I acknowledge the duty to disclose information which is material to the examination of this application in accordance with Title 37, Code of Federal Regulations, §1.56(a).

Japanese Language Declaration

私は、合衆国法典第35部第119条にもとづく下記の外国特許出願または発明者証出願の外国優先権利益を主張し、さらに優先権の主張に係わる基礎出願の出願日前の出願日を有する外国特許出願または発明者証出願を以下に明記する：

Prior foreign applications

先の外国出願

11-305026	Japan	27/10/1999
(Number) (番号)	(Country) (国名)	(Day/Month/Year Filed) (出願の年月日)
_____ (番号)	_____ (国名)	_____ (出願の年月日)
_____ (番号)	_____ (国名)	_____ (出願の年月日)

Priority claimed

優先権の主張

<input checked="" type="checkbox"/> Yes あり	<input type="checkbox"/> No なし
<input type="checkbox"/> Yes あり	<input type="checkbox"/> No なし
<input type="checkbox"/> Yes あり	<input type="checkbox"/> No なし

私は、合衆国法典第35部第120条にもとづく下記の合衆国特許出願の利益を主張し、本願の請求の範囲各項に記載の主題が合衆国法典第35部第112条第1項に規定の様様で先の合衆国出願に開示されていない限度において、先の出願の出願日と本願の国内出願日またはPCT国際出願日の間に公表された連邦規則法典第37部第1章第56条(a)項に記載の所要の情報を開示すべき義務を有することを認める：

0/	
(Application Serial No.) (出願番号)	(Filing Date) (出願日)
0/	
(Application Serial No.) (出願番号)	(Filing Date) (出願日)

I hereby claim the benefit under Title 35, United States Code, §120 of any United States application(s) listed below and, insofar as the subject matter of each of the claims of this application is not disclosed in the prior United States application in the manner provided by the first paragraph of Title 35, United States Code, §112, I acknowledge the duty to disclose material information as defined in Title 37, Code of Federal Regulations, §1.56(a) which occurred between the filing date of the prior application and the national or PCT international filing date of this application:

(現況) (特許済み、係属中、放棄済み)	(Status) (patented, pending, abandoned)
(現況) (特許済み、係属中、放棄済み)	(Status) (patented, pending, abandoned)

私は、ここに自己の知識にもとづいて行った陳述がすべて真実であり、自己の有する情報および信ずるところに従って行った陳述が真実であると信じ、さらに故意に虚偽の陳述等を行った場合、合衆国法典第18部第1001条により、罰金もしくは禁錮に処せられるか、またはこれらの刑が併科され、またかかる故意による虚偽の陳述が本願ないし本願に対して付与される特許の有効性を損うことがあることを認識して、以上の陳述を行ったことを宣言する。

I hereby declare that all statements made herein of my own knowledge are true and that all statements made on information and belief are believed to be true; and further that these statements were made with the knowledge that willful false statements and the like so made are punishable by fine or imprisonment, or both, under Section 1001 of Title 18 of the United States Code and that such willful false statements may jeopardize the validity of the application or any patent issued thereon.

Under the Paperwork Reduction Act of 1995, no persons are required to respond to a collection of information unless it displays a valid OMB control number.

Japanese Language Declaration
(日本語宣言書)

委任状: 私は下記の発明者として、本出願に関する一切の手続きを米特許商標局に対して遂行する弁理士または代理人として、下記の者を指名いたします。(弁理士、または代理人の氏名及び登録番号を明記のこと)

POWER OF ATTORNEY: As a named inventor, I hereby appoint the following attorney(s) and/or agent(s) to prosecute this application and transact all business in the Patent and Trademark Office connected therewith (list name and registration number)

Attorney	Reg. No.
Patrick G. Burns	29,367
Roger D. Greer	26,174
Lawrence J. Crain	31,497
Steven P. Fallon	35,132

Attorney	Reg. No.
James K. Folker	37,538
Jonathan D. Feuchtwang	41,017
B. Joe Kim	41,895
Joel H. Bootzin	42,343

直接電話連絡先: (名前及び電話番号)

Send Correspondence to:

Direct Telephone Calls to: (name and telephone number)

Patrick G. Burns, Esq.
Greer, Burns & Crain, Ltd.
Sears Tower - Suite 8660
Chicago, IL 60606 (312) 993-0080

唯一または第一発明者名	Full name of sole or first inventor
	Masatoshi HARAGUCHI
発明者の署名	Inventor's signature
日付	Date
	Masatoshi Haraguchi June 10, 2000
住所	Residence
	Kawasaki-shi, Kanagawa, Japan
国籍	Citizenship
	Japanese
私書箱	Post Office Address
	c/o FUJITSU LIMITED, 1-1, Kamikodanaka 4-chome, Nakahara-ku, Kawasaki-shi, Kanagawa 211-8588 JAPAN
第二共同発明者	Full name of second joint inventor, if any
	Masakazu HAYASHI
第二共同発明者	Second inventor's signature
日付	Date
	Masakazu Hayashi June 10, 2000
住所	Residence
	Kawasaki-shi, Kanagawa, Japan
国籍	Citizenship
	Japanese
私書箱	Post Office Address
	c/o FUJITSU LIMITED, 1-1, Kamikodanaka 4-chome, Nakahara-ku, Kawasaki-shi, Kanagawa 211-8588 JAPAN

(第三以降の共同発明者についても同様に記載し、署名をすること)

(Supply similar information and signature for third and subsequent joint inventors.)

第三共同発明者名		Full name of third joint inventor, if any Yuji WATANABE	
第三共同発明者の署名	日付	Third inventor's signature <i>Yuji Watanabe</i>	Date June 10, 2000
住所	Residence Kawasaki-shi, Kanagawa, Japan		
国籍	Citizenship Japanese		
私書箱	Post Office Address c/o FUJITSU LIMITED, 1-1, Kamikodanaka 4-chome, Nakahara-ku, Kawasaki-shi, Kanagawa 211-8588 JAPAN		
第四共同発明者名		Full name of fourth joint inventor, if any	
第四共同発明者の署名	日付	Fourth inventor's signature	Date
住所	Residence		
国籍	Citizenship		
私書箱	Post Office Address		
第五共同発明者名		Full name of fifth joint inventor, if any	
第五共同発明者の署名	日付	Fifth inventor's signature	Date
住所	Residence		
国籍	Citizenship		
私書箱	Post Office Address		
第六共同発明者名		Full name of sixth joint inventor, if any	
第六共同発明者の署名	日付	Sixth inventor's signature	Date
住所	Residence		
国籍	Citizenship		
私書箱	Post Office Address		